

Tiling Periodicity

Yury Lifshits
joint work with Juhani Karhumäki

Steklov Institute of Mathematics, St.Petersburg, Russia
yura@logic.pdmi.ras.ru

Dagstuhl
May 2006

Classical Notion of Periodicity

The string S is called **purely periodic** if

$$S = W^k = W \dots W$$

Equivalently

$$\forall 1 \leq i < i + p \leq n : s_i = s_{i+p}$$

Classical Notion of Periodicity

The string S is called **purely periodic** if

$$S = W^k = W \dots W$$

Equivalently

$$\forall 1 \leq i < i+p \leq n: s_i = s_{i+p}$$

Is the following string purely periodic?

| | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| A | A | B | B | A | A | B | B | C | C | D | D | C | C | D | D |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|

Classical Notion of Periodicity

The string S is called **purely periodic** if

$$S = W^k = W \dots W$$

Equivalently

$$\forall 1 \leq i < i + p \leq n : s_i = s_{i+p}$$

Is the following string purely periodic?

| | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| A | A | B | B | A | A | B | B | C | C | D | D | C | C | D | D |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|

Not in the classical sense. But...

Outline of the Talk

- 1 Notion of Tiling Periodicity
- 2 Minimal Tiling Period Conjecture
- 3 Properties of Tiling Periodicity
 - Maximal Number of Periods
 - Relation to Classical Periodicity
 - Algorithm for Finding Minimal Tiling Periods
- 4 Future Work

Motivating Examples

A A B B A A B B C C D D C C D D

The string above is not periodic, but pink **structure**

A A B B A A B B C C D D C C D D

is a kind of period, since we can cover initial string by **four parallel copies** of it:

A A B B A A B B C C D D C C D D

Motivating Examples

A A B B A A B B C C D D C C D D

The string above is not periodic, but pink **structure**

A A B B A A B B C C D D C C D D

is a kind of period, since we can cover initial string by **four parallel copies** of it:

A A B B A A B B C C D D C C D D

The simplest example:

A A B B

Formal Definition

A **tiling string** (or tiler) is a string over $\Sigma \cup \square$ alphabet, where \square is a special **transparent** (or undefined) letter. Sometimes the term **partially defined word** is also used

Formal Definition

A **tiling string** (or tiler) is a string over $\Sigma \cup \square$ alphabet, where \square is a special **transparent** (or undefined) letter. Sometimes the term **partially defined word** is also used

A tiling string S is called the **tiling period** of (ordinary) string T if we can cover T by parallel copies of S satisfying the following:

- All defined (visible) letters of S -copies match the text letters
- Every text letter covered by **exactly one** defined (visible) letter

Why Tiling Periodicity

- New structural properties of texts (Conjecture: tiling periodicity is not expressible in word equations)

Why Tiling Periodicity

- New structural properties of texts (Conjecture: tiling periodicity is not expressible in word equations)
- New tool for text compression

Why Tiling Periodicity

- New structural properties of texts (Conjecture: tiling periodicity is not expressible in word equations)
- New tool for text compression
- Relations to multidimensional periodicity

Why Tiling Periodicity

- New structural properties of texts (Conjecture: tiling periodicity is not expressible in word equations)
- New tool for text compression
- Relations to multidimensional periodicity
- Natural generalization of the classical notion

Why Tiling Periodicity

- New structural properties of texts (Conjecture: tiling periodicity is not expressible in word equations)
- New tool for text compression
- Relations to multidimensional periodicity
- Natural generalization of the classical notion
- Pattern discovery (?????)

Partial Order on Tilers

We say that one tiling string (tiler) S is **smaller** than another tiler Q , if Q can be covered by several parallel copies of S satisfying the following:

- All defined (visible) letters of S -copies match the visible Q letters
- Every Q letter covered by **exactly one** defined (visible) letter

Example:

| | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| A | A | B | B | A | A | B | B | C | C | D | D | C | C | D | D |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|

is less than

| | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| A | A | B | B | A | A | B | B | C | C | D | D | C | C | D | D |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|

Minimal Tiling Period Conjecture

Main Conjecture: For every ordinary string there exists a unique minimal tiling period (it is less than any other tiling period).

Reformulation Any two tiling periods have a common tiling “subperiod”

Minimal Tiling Period Conjecture

Main Conjecture: For every ordinary string there exists a unique minimal tiling period (it is less than any other tiling period).

Reformulation Any two tiling periods have a common tiling “subperiod”

Big surprise (at least for me): conjecture is wrong! Look at (minimal known) counterexample:

| | | | | | | | | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| A | A | A | A | A | A | A | A | B | A | A | B | B | A | A | B | A | A | A | A | A | A | A | A |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|

and

| | | | | | | | | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| A | A | A | A | A | A | A | A | B | A | A | B | B | A | A | B | A | A | A | A | A | A | A | A |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|

How Many Tiling Periods?

Let $L(n)$ be the number of periods of the string of length n over a unary alphabet. Then:

- $L(1)=1$
- For every $n > 1$ we can compute L by recursive formula:

$$L(1) = 2; L(n) = \sum_{d|n, d \neq n} L(d)$$

- $L(36) = 52$
- $L(p_1 \cdots p_k) = (k + 1)!$
- (To be done) What is the upper limit of $L(n)/n$?

Tiling Periods are Always Smaller than Classical

Theorem Take any pair of tiling period and classical period. Then they have a common “tiling subperiod”. Any minimal tiling period of string T is also a tiling period of any classical period of T .

Finding Minimal Tiling Periods: Sketch

- 1 Define a notion of “ranged periodicity”
- 2 Prove that any minimal tiling root corresponds to the “best” chain of embedded ranged periodicities
- 3 Find all ranged periodicities
- 4 Find the “best” chain

Directions for Further Research

- Study **not pure** tiling periodicity
- How often strings are tiling periodic?
- Whether the property **“string has a tiling square root”** can be expressed by word equations?
- Whether all minimal tiling roots have the same number of visible letters?
- Find natural sources of tiling periodicity
- Improve the complexity of the algorithm for finding minimal tiling periods
- Find relevant references

Summary

Main points:

- New notion: tiling periodicity



Summary

Main points:

- New notion: tiling periodicity



- The minimal tiling root is not necessary unique!

Summary

Main points:

- New notion: tiling periodicity



- The minimal tiling root is not necessary unique!
- Algorithm for finding minimal tiling roots

Last Slide

Related paper will appear soon at
<http://logic.pdmi.ras.ru/~yura/>

Last Slide

Related paper will appear soon at
<http://logic.pdmi.ras.ru/~yura/>

Thanks for inviting to the seminar
Thanks for opportunity for the second talk
Thanks for attention

Last Slide

Related paper will appear soon at
<http://logic.pdmi.ras.ru/~yura/>

Thanks for inviting to the seminar
Thanks for opportunity for the second talk
Thanks for attention

Questions?